風水
Heap Feng Shui in JavaScript

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• What is Heap Feng Shui?
  ○ the ancient art of arranging heap blocks in order to redirect the program control flow to the shellcode

• Heap Feng Shui in JavaScript
  ○ precise application data overwrites
  ○ reliable browser exploitation
Overview

- State of the art in browser exploitation
- Internet Explorer heap internals
- HeapLib JavaScript library
- Heap manipulation
- Mitigation
Part I

State of the art in browser exploitation
Stack overflows

Very hard to exploit in most cases:

<table>
<thead>
<tr>
<th>Target</th>
<th>Protection</th>
</tr>
</thead>
<tbody>
<tr>
<td>return address</td>
<td>stack cookies (/GS flag)</td>
</tr>
<tr>
<td>SEH frame</td>
<td>SafeSEH exception handler table</td>
</tr>
<tr>
<td>local variables</td>
<td>local variable reordering in the Visual C++ compiler</td>
</tr>
</tbody>
</table>
Heap overflows

Generic heap exploitation is also difficult:

<table>
<thead>
<tr>
<th>Target</th>
<th>Protection</th>
</tr>
</thead>
<tbody>
<tr>
<td>doubly-linked list of free chunks</td>
<td>safe unlinking</td>
</tr>
<tr>
<td>heap chunk header</td>
<td>8-bit header cookie in XP, XOR of the header data in Vista</td>
</tr>
<tr>
<td>lookaside linked list</td>
<td>removed in Vista</td>
</tr>
</tbody>
</table>
What's left?

- Non-array stack overflows
  - very rare
- Use of uninitialized variables
  - stack variables
  - use after free
- Application data on the heap
  - application specific memory allocators
  - function pointers
  - C++ object pointers
WebView setSlice exploit

- Uses heap spraying to fill the browser heap with shellcode
- Overwrites application data in the previous heap chunk
- Multiple attempts until it either hits an object pointer, or crashes
Heap spraying

Developed by Blazde and SkyLined, used by most browser exploits since 2004.

```javascript
var x = new Array();

// Fill 200MB of memory with copies of the NOP slide and shellcode
for (var i = 0; i < 200; i++) {
    x[i] = nop + shellcode;
}
```
Normal heap layout

used memory:
free memory:
After heap spraying

used memory:  
free memory:  
shellcode:  

Address 0x0C0C0C0C0C is very likely to contain shellcode.
Function pointer overwrite

1. Spray the heap with 200MB of shellcode
2. Overwrite a function pointer with \texttt{0xC0C0C0C0C0}
3. Call the function pointer
1. Spray the heap with 200MB of shellcode, using byte 0xC as a nop slide

2. Overwrite an object pointer with 0x0C0C0C0C0C

3. Call a virtual function of the object
Unreliable exploitation

- Heap spraying is a great technique, but the setSlice exploit is still not reliable
- Overwriting application data requires a specific layout of heap chunks
- We need to control the heap state
Part II

Heap Feng Shui
Heap Feng Shui

• The heap allocator is deterministic

• Specific sequences of allocations and frees can be used to control the layout

used: □
free: □
Heap Feng Shui

- The heap allocator is deterministic
- Specific sequences of allocations and frees can be used to control the layout

![Diagram showing allocation and free space]

used:  
free:  
our data: 

We allocate two 4KB blocks
• The heap allocator is deterministic

• Specific sequences of allocations and frees can be used to control the layout

used: We free the first 4KB block
free:
our data:
Heap Feng Shui

- The heap allocator is deterministic
- Specific sequences of allocations and frees can be used to control the layout

The application allocates a 4KB block and reuses our data
Heap Feng Shui

• The heap allocator is deterministic

• Specific sequences of allocations and frees can be used to control the layout

We just exploited an uninitialized data vulnerability

used:  free:  our data:  We just exploited an uninitialized data vulnerability
Heap Feng Shui in JavaScript

- We want to set the heap state before triggering a vulnerability
- Heap spraying proves that JavaScript can access the system heap
- We need a way to allocate and free blocks of an arbitrary size
Part III

Internet Explorer heap internals
Internet Explorer heap usage

- JavaScript runtime
  - objects
  - JavaScript heap
- MSHTML engine
  - strings
  - Default process heap
- ActiveX objects
  - Dedicated heaps
JavaScript strings

The string "AAAA" is stored as:

<table>
<thead>
<tr>
<th>string size</th>
<th>string data</th>
<th>null terminator</th>
</tr>
</thead>
<tbody>
<tr>
<td>4 bytes</td>
<td>length / 2 bytes</td>
<td>2 bytes</td>
</tr>
<tr>
<td>08 00 00 00</td>
<td>41 00 41 00 41 00</td>
<td>00 00</td>
</tr>
</tbody>
</table>

We can calculate its size in bytes with:

\[
\text{bytes} = \text{len} \times 2 + 6 \\
\text{len} = (\text{bytes} - 6) / 2
\]
var str1 = "AAAAAAAAAAA"; // no allocation

// allocates a 10 character string
var str2 = str1.substr(0, 10);

// allocates a 20 character string
var str3 = str1 + str2;
String garbage collection

- Mark-and-sweep algorithm, frees all unreferenced objects
- Triggered by a number of heuristics
- Explicitly by the `CollectGarbage()` call in Internet Explorer
var padding = "AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA..."
var str;

function alloc(bytes) {
    str = padding.substr(0, (bytes-6)/2);
}

function free() {
    str = null;
    CollectGarbage();
}

alloc(0x10000); // allocate 64KB memory block
free(); // free memory block
OLEAUST32 allocator

Not all string allocations and frees reach the system memory allocator

• custom memory allocator in OLEAUST32
• caching of free memory blocks
• 4 bins for blocks of different sizes
• up to 6 free blocks stored in each bin
bin = the right bin for the requested size

if (bin not empty)
    find a block in the bin > requested size
    if (found)
        return block
    else
        return sysalloc(size)
else
    return sysalloc(size)
OLEAUT32 free function

bin = the right bin for the block size

if (bin not full)
    add block to bin
else
    find the smallest block in the bin
    if (smallest block < new block)
        sysfree(smallest block)
        add new block to bin
    else
        sysfree(new block)
Bypassing the cache

- Our freed blocks will go into the cache
- Freeing 6 maximum sized blocks for each bin will push all smaller blocks out
- Allocating the 6 blocks again will leave the cache empty
- When the cache is empty, allocations will come from the system heap
Plunger Technique

1. Allocate 6 maximum size blocks
2. Allocate our blocks
3. Free our blocks
4. Free 6 maximum size blocks
5. Allocate 6 maximum size blocks

OLEAUT32 cache

maximum size blocks

empty
Plunger Technique

1. Allocate 6 maximum size blocks
2. Allocate our blocks
3. Free our blocks
4. Free 6 maximum size blocks
5. Allocate 6 maximum size blocks
Plunger Technique

1. Allocate 6 maximum size blocks
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maximum size blocks

our blocks

OLEAUT32 cache
Plunger Technique

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Part IV

HeapLib - JavaScript heap manipulation library
Introducing HeapLib

- Supports Internet Explorer 5-7
- Object oriented API
- Functions for:
  - heap logging and debugging
  - allocation and freeing of blocks with arbitrary size and contents
  - high-level heap manipulation function (not yet supported on Vista)
<script src="heapLib.js"></script>

<script>
    var heap = new heapLib.ie();
    heap.gc();

    heap.debugHeap(true);

    heap.alloc(512);
    heap.alloc("BBBBBB", "foo");

    heap.free("foo");

    heap.debugHeap(false);
</script>
HeapLib Demo
Part V

Windows Heap Manipulation
Windows Heap Overview

Pre-Vista

Heap

FreeList[0] → 1024 → 2080 → 8192
FreeList[1] → 8
...
FreeList[127] → 1016 → 1016

Lookaside

Lookaside Table

Lookaside[0] → 8
Lookaside[1] → 8
...
Lookaside[126] → 1016
Free Algorithm

if size >= 512KB
    free with VirtualFree
    return

if size < 1KB and lookaside not full
    add to lookaside list
    return

coalesce block with free blocks around it

if size < 1KB
    add to FreeList[size/8]
else
    add to FreeList[0]
Allocate Algorithm

if size >= 512KB
    alloc with VirtualAlloc
    return

if size < 1KB
    if lookaside not empty
        return a block from the lookaside
    if FreeList[size/8] not empty
        return a block from FreeList[size/8]

if FreeList[0] not empty
    return a block from FreeList[0]

allocate more memory with VirtualAlloc
Defragmenting the heap

To allocate two consecutive blocks, we need to defragment the heap.

```javascript
for (var i = 0; i < 1000; i++)
    heap.alloc(0x2010);
```

used:  □
free: □
Defragmenting the heap

To allocate two consecutive blocks, we need to defragment the heap.

```javascript
for (var i = 0; i < 1000; i++)
    heap.alloc(0x2010);
```

![Diagram showing defragmentation process]
Putting a block on the FreeList

To put a block on the free list, we need to ensure that it is not coalesced.

```c
heap.alloc(0x2010, "foo");
heap.alloc(0x2010);
heap.alloc(0x2010, "foo");
heap.free("foo");
```
Putting a block on the FreeList

To put a block on the free list, we need to ensure that it is not coalesced.

```c
heap.alloc(0x2010, "foo");
heap.alloc(0x2010);
heap.alloc(0x2010, "foo");
heap.free("foo");
```

used:  
free:  
our blocks:  

↓

```
Emptying the lookaside

To empty the lookaside, allocate enough blocks of the same size.

```javascript
for (var i = 0; i < 100; i++)
    heap.alloc(512);
```
Freeing to the lookaside

To put a block on the lookaside, empty it and free the block.

```javascript
for (var i = 0; i < 100; i++)
    heap.alloc(512);

heap.alloc(512, "foo");
heap.free("foo");
```
Object pointer overwrite

The lookaside linked list can be used to exploit object pointer overwrites without heap spraying.

1. Empty the lookaside
2. Build a fake vtable block
3. Free the fake vtable to the lookaside
4. Overwrite an object pointer with the address of the lookaside head
5. Call a virtual function of the object
Object pointer overwrite

```assembly
mov ecx, dword ptr [eax] ; get the vtable address
push eax                ; push the 'this' pointer
call dword ptr [ecx+08h] ; call virtual func
```

NULL disassembles as two `sub [eax], al` instructions
Exploit Demo
Mitigation

- Heap isolation
- Non-determinism in the heap allocator
Questions?

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